# ON LOCAL IRREGULARITY OF THE VERTEX COLORING OF THE CORONA PRODUCT OF A TREE GRAPH 

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#### Abstract

Let $G=(V, E)$ be a graph with a vertex set $V$ and an edge set $E$. The graph $G$ is said to be with a local irregular vertex coloring if there is a function $f$ called a local irregularity vertex coloring with the properties: (i) $l:(V(G)) \rightarrow\{1,2, \ldots, k\}$ as a vertex irregular $k$-labeling and $w: V(G) \rightarrow N$, for every $u v \in E(G)$, $w(u) \neq w(v)$ where $w(u)=\sum_{v \in N(u)} l(i)$ and (ii) $\operatorname{opt}(l)=\min \left\{\max \left\{l_{i}: l_{i}\right.\right.$ is a vertex irregular labeling $\left.\}\right\}$. The chromatic number of the local irregularity vertex coloring of $G$ denoted by $\chi_{l i s}(G)$, is the minimum cardinality of the largest label over all such local irregularity vertex colorings. In this paper, we study a local irregular vertex coloring of $P_{m} \bigodot G$ when $G$ is a family of tree graphs, centipede $C_{n}$, double star graph ( $S_{2, n}$ ), Weed graph $\left(S_{3, n}\right)$, and $E$ graph $\left(E_{3, n}\right)$.


Keywords: Local irregularity, Corona product, Tree graph family.

## 1. Introduction

Let $G(V, E)$ be a connected and simple graph with a vertex set $V$ and an edge set $E$. In this paper, we combine two concepts, namely the local antimagic vertex coloring and the distance irregular labelling, with a local irregularity of vertex coloring. This concept firstly was introduced by Kristiana [2, 3], et. al. The latest research was conducted by Azzahra [4], who examined the local irregularity vertex coloring of a grid graph family. In this paper we study the local irregularity of vertex coloring of corona product graph of a tree graph family.

Definition 1. Suppose $l: V(G) \rightarrow\{1,2, \ldots, k\}$ and $w: V(G) \rightarrow N$, where

$$
w(u)=\sum_{v \in N(u)} l(v)
$$

then $l(v)$ is called the vertex irregular $k$-labeling and $w(u)$ is called the local irregularity of vertex coloring if
(i) $\operatorname{opt}(l)=\min \left\{\max \left\{l_{i}\right\}: l_{i}\right.$ vertex irregular labeling $\}$;
(ii) for every $u v \in E(G), w(u) \neq w(v)$.

Definition 2. The chromatic number of local irregular graph $G$ denoted by $\chi_{l i s}(G)$, is the minimum of cardinality of the local irregularity of vertex coloring.

In this paper, we will use the following lemma which gives a lower bound on the chromatic number of local irregular vertex coloring:

Lemma 1 [2]. Let $G$ be a simple and connected graph, then $\chi_{l i s}(G) \geq \chi(G)$.
Proposition 1 [2]. Let $G$ be a graph each two adjacent vertices of which have a different vertex degree then $\operatorname{opt}(l)=1$.

Proposition 2 [2]. Let $G$ be a graph each two adjacent vertices have the same vertex degree then $\operatorname{opt}(l) \geq 2$.

Definition 3 [1]. Let $G$ and $H$ be two connected graphs. Let o be a vertex of H. The corona product of the combination of two graphs $G$ and $H$ is defined as the graph obtained by taking a duplicate of graph $G$ and $|V(G)|$ a duplicate of graph $H$, namely $H_{i} ; i=1,2,3, \ldots|V(G)|$ then connects each vertex $i$ in $G$ to each vertex in $H_{i}$. The corona product of the graphs $G$ and $H$ is denoted by $G \odot H$.

## 2. Result and discussion

In this paper, we analyze the new result of the chromatic number of local irregular vertex coloring of corona product by family of tree graph $\left(P_{m} \odot G\right)$ where $G$ is centipede graph $\left(C_{n}\right)$, double star graph $\left(S_{2, n}\right)$, and Weed graph $\left(S_{3, n}\right)$.

Theorem 1. Let $G=P_{m} \odot C p_{n}$, be a corona product of a path graph of order $m$ and a centipede graph of order $n$ for $n, m \geq 2$, then

$$
\chi_{l i s}\left(P_{m} \bigodot C p_{n}\right)=\left\{\begin{array}{l}
5, \text { for } m=3 \text { and } n=2,3, \\
6, \\
\text { for } m=2 \text { and } n=2,3 \text { or for } m=3 \text { and } n \geq 4, \\
7, \text { for } m=2 \text { and } n \geq 4 \\
8, \\
\text { for } m \geq 4
\end{array} \text { or for } m \geq 4 \text { and } n=2,3,\right.
$$

with $\operatorname{opt}(l)$ defined as

$$
\operatorname{opt}(l)\left(P_{m} \bigodot C p_{n}\right)=\left\{\begin{array}{llll}
1, & \text { for } m=3 & \text { and } n=3, \\
1,2, & \text { for } m=2 & \text { and } n=2 & \text { or } \\
& \text { for } m=3 & \text { and } n=2 & \text { or } \\
& \text { for } m \geq 3 & \text { and } n \geq 4 .
\end{array}\right.
$$

Proof. Vertex set is

$$
V\left(P_{m} \bigodot C P_{n}\right)=\left\{x_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{y_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\}
$$

and the edge set is

$$
\begin{gathered}
E\left(P_{m} \bigodot C P_{n}\right)=\left\{x_{i} x_{i+1} ; 1 \leq i \leq m-1\right\} \cup\left\{x_{i j} x_{i j+1} ; 1 \leq i \leq m, 1 \leq j \leq n-1\right\} \\
\cup\left\{x_{i j} y_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{x_{i} x_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{x_{i} y_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\},
\end{gathered}
$$

the order and size respectively are $2 m n+m$ and $4 m n-1$.
Case 1: $m \neq p, m \geq 2, p \geq 2, n \geq 3$.

First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=4$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=4$, if $l\left(x_{1}\right)=l\left(x_{3}\right)=1, l\left(x_{2}\right)=2$, $l\left(x_{i j}\right)=l\left(y_{i j}\right)=1$ then $w\left(x_{1}\right)=w\left(x_{2}\right)$, then there are 2 adjacent vertices that have the same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \quad 1 \leq i \leq 3, \quad j=1, \quad l\left(y_{i j}\right)=2, \quad 1 \leq i \leq 3, \quad j=2, \\
l\left(x_{i}\right)=1 \rightarrow w\left(x_{i}\right) \neq w\left(x_{i+1}\right), \quad w\left(x_{i 1}\right) \neq\left(x_{i 2}\right),
\end{gathered}
$$

then $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 5$. Based on this, we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 5$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$. Furthermore, the upper bound for the chromatic number of local irregular $\left(P_{m} \odot C p_{n}\right)$, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(x_{i j}\right)=1, \\
l\left(y_{i j}\right)=\left\{\begin{array}{lll}
1, & \text { for } \quad 1 \leq i \leq 3 \text { and } j=1, \\
2, & \text { for } 1 \leq i \leq 3 \text { and } j=2 .
\end{array}\right.
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}6, & \text { for } i=1,3, \\
8, & \text { for } i=2,\end{cases} \\
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } 1 \leq i \leq 3 \text { and } j=1, \\
4, & \text { for } 1 \leq i \leq 3 \text { and } j=2,\end{cases} \\
w\left(y_{i j}\right)=2, \quad \text { for } 1 \leq i \leq 3 \text { and } j=1,2 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 5$, and we have $5 \leq \chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 5$, so $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$ for $m=3$ and $n=2$.

Case 2: $m=n=3$.
Based on Proposition 1, $\operatorname{opt}(l)=1$. So the lower bound of $\left(P_{m} \odot C p_{n}\right)$ is

$$
\chi_{l i s}\left(P_{m} \bigodot C p_{n}\right) \geq 5
$$

Hence $\operatorname{opt}(l)=1$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}7, & \text { for } i=1,3, \\
8, & \text { for } i=2,\end{cases} \\
w\left(y_{i j}\right)= \begin{cases}3, & \text { for } 1 \leq i \leq 3 \text { and } j \equiv 1,3(\bmod 4), \\
4, & \text { for } 1 \leq i \leq 3 \text { and } j=2,\end{cases} \\
w\left(x_{i j}\right)=2, \quad \text { for } 1 \leq i \leq 3 \text { and } 1 \leq j \leq 3 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 5$. We have $5 \leq \chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 5$, so $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$ for $m=3$ and $n=3$.

Case 3: $m=n=2$.
First step here is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, if $l\left(x_{1}\right)=1, l\left(x_{2}\right)=2, l\left(x_{i j}\right)=l\left(y_{i j}\right)=1$, then $w\left(x_{11}\right)=w\left(x_{12}\right)$ and there are 2 adjacent vertices, that have the same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{1}\right)=1, \quad l\left(x_{2}\right)=2, \quad l\left(x_{i j}\right)=1, \quad l\left(y_{i 1}\right)=1, \quad i=1,2, \quad l\left(y_{i 2}\right)=2, \quad i=1,2,
$$

then $w\left(x_{1}\right) \neq w\left(x_{2}\right), w\left(x_{i 1}\right) \neq w\left(x_{i 2}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 6$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
l\left(x_{i}\right)=\left\{\begin{array}{ll}
1, & \text { for } i=1, \\
2, & \text { for } i=2,
\end{array} \quad l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)= \begin{cases}1, & \text { for } i=1,2 \text { and } j=1, \\
2, & \text { for } i=1,2 \text { and } j=2 .\end{cases}\right.
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}6, & \text { for } i=2, \\
7, & \text { for } i=1,\end{cases} \\
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i=1 \text { and } j=1 \\
4, & \text { for } i=1 \text { and } j=1, \text { or for } i=2 \text { and } j=1, \\
5, & \text { for } i=2 \text { and } j=2,\end{cases} \\
w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i=1 \text { and } j=1,2, \\
3, & \text { for } i=2 \text { and } j=1,2 .\end{cases}
\end{gathered}
$$

We have the following upper bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 6$. We have $6 \leq \chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 6$, so $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$ for $m=2$ and $n=2$.

Case 4: $m=2$ and $n=3$.
First step here is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, if

$$
l\left(x_{i}\right)=l\left(x_{i j}\right)=1, \quad l\left(y_{1 j}\right)=1, \quad l\left(y_{2 j}\right)=1, \quad j=1,2, \quad l\left(y_{i 3}\right)=2,
$$

then $w\left(x_{22}\right)=w\left(x_{23}\right)$, so there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{1}\right)=1, \quad l\left(x_{2}\right)=2, \quad l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1,
$$

then $w\left(x_{1}\right) \neq w\left(x_{2}\right), w\left(x_{i, 1}\right) \neq w\left(x_{i, 2}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 6$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
l\left(x_{i}\right)=\left\{\begin{array}{ll}
1, & \text { for } i=1, \\
2, & \text { for } i=2,
\end{array} \quad l\left(x_{i j}\right)=1, \quad w\left(y_{i j}\right)=1 .\right.
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}7, & \text { for } i=2, \\
8, & \text { for } i=1,\end{cases} \\
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i=1 \text { and } j=1,3, \\
4, & \text { for } i=1 \text { and } j=2, \text { or for } i=2 \text { and } j=1,3, \\
5, & \text { for } i=2 \text { and } j=2,\end{cases} \\
w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i=1 \text { and } 1 \leq j \leq 3, \\
3, & \text { for } i=2 \text { and } 1 \leq j \leq 3 .\end{cases}
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 6$. So we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$ for $m=2$ and $n=3$.

Case 5: $m=3$ and $n \geq 4$.
First step to prove this theorem in this case is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, if $l\left(x_{1}\right)=l\left(x_{3}\right)=1, l\left(x_{2}\right)=2, l\left(x_{i j}\right)=l\left(y_{i j}\right)=1$, then $w\left(x_{1}\right)=$ $w\left(x_{2}\right)$ so there are 2 adjacent vertices with the have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \quad 1 \leq i \leq 3, \quad j=1, n, \quad j \equiv 0(\bmod 2), \\
l\left(y_{i j}\right)=2, \quad 1 \leq i \leq 3, \quad j \equiv 1,3(\bmod 4), \quad j \neq 1, n,
\end{gathered}
$$

with the $w\left(x_{i}\right) \neq w\left(x_{i+1}\right), w\left(x_{i j}\right)=w\left(x_{i j+1}\right)$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 6$.

After that, we will find the upper bound for $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
l\left(y_{i j}\right)=\left\{\begin{array}{l}
1, \quad \text { for } 1 \leq i \leq 3 \text { and } j=1, n \text { or for } 1 \leq i \leq 3 \text { and } j \equiv 0(\bmod 2), \\
2, \quad \text { for } 1 \leq i \leq 3 \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n .
\end{array}\right.
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}2 n+n / 2, & \text { for } i=1,3 \text { and } n \equiv 0(\bmod 2), \\
2 n+\lfloor n / 2\rfloor, & \text { for } i=1,3 \text { and } n \equiv 1,3(\bmod 4), \\
3 n+1-n / 2, & \text { for } i=2 \text { and } n \equiv 0(\bmod 2), \\
3 n+1-\lceil n / 2\rceil, & \text { for } i=2 \text { and } n \equiv 1,3(\bmod 4),\end{cases} \\
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } 1 \leq i \leq 3 \text { and } j=1, n, \\
4, & \text { for } 1 \leq i \leq 3 \text { and } j \equiv 0(\bmod 2), \\
5, & \text { for } 1 \leq i \leq 3 \text { and } j \equiv 1,3(\bmod 4), \quad j \neq 1, n, \\
w\left(y_{i j}\right)=2 .\end{cases}
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 6$. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$ for $m=3$ and $n \geq 4$.
Case 6: $m \equiv 0(\bmod 2), m \geq 4$ and $n=2$.

First step here is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, if

$$
l\left(x_{i}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \quad l\left(x_{i}\right)=2, \quad i \equiv 0(\bmod 4), \quad l\left(x_{i j}\right)=l\left(y_{i j}\right)=1
$$

then $w\left(x_{i j}\right)=w\left(x_{i j+1}\right)$, then there are 2 adjacent vertices that have same color, this contradicts the definition of vertex coloring. If

$$
\begin{array}{cl}
l\left(x_{i}\right)=1, & i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \quad l\left(x_{i}\right)=2, \quad i \equiv 0(\bmod 4) \\
& l\left(x_{i j}\right)=1 l\left(y_{i j}\right)=1, \quad j=2, \quad l\left(y_{i j}\right)=2, \quad j=1,
\end{array}
$$

then $w\left(x_{i j}\right) \neq w\left(x_{i j+1}\right) ; w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right)$. So we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 7$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{aligned}
& l\left(x_{i}\right)= \begin{cases}1, & \text { for } i \equiv 1,3(\bmod 4) \text { or for } i \equiv 2(\bmod 4) \\
2, & \text { for } i \equiv 0(\bmod 4),\end{cases} \\
& l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)= \begin{cases}1, & \text { for } 1 \leq i \leq m \text { and } j=1 \\
2, & \text { for } 1 \leq i \leq m \text { and } j=2\end{cases}
\end{aligned}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}6, & \text { for } i=1, m, \\
7, & \text { for } i \equiv 0(\bmod 2), \quad i \neq m, \\
8, & \text { for } i \equiv 1,3(\bmod 4), \quad i \neq 1,\end{cases} \\
w\left(x_{i j}\right)=\left\{\begin{array}{lr}
3, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=1 \quad \text { or } \quad \text { for } i \equiv 2(\bmod 4) \text { and } j=1, \\
4, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=2 \quad \text { or } \quad \text { for } i \equiv 2(\bmod 4) \text { and } j=2 \text { or } \\
5, & \text { for } i \equiv 0(\bmod 4) \text { and } j=1,
\end{array}\right. \\
w\left(y_{i j}\right)=\left\{\begin{array}{ll}
2, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=2, \\
3, & \text { for } i=0(\bmod 4) \text { and } j=1,2 .
\end{array} \quad \text { or } \quad \text { for } i \equiv 2(\bmod 4) \text { and } j=1,2,\right.
\end{gathered}
$$

We have the upper bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$ for $m \geq 4$ and $n=2$.
Case 7: $m \equiv 0(\bmod 2), m \geq 4$ and $n=3$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=5$, in this case if

$$
l\left(x_{i}\right)=l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \quad 1 \leq i \leq m, \quad j=3, \quad l\left(y_{i j}\right)=2, \quad 1 \leq i \leq m, \quad j=1,2,
$$

then $w\left(x_{i}\right)=w\left(x_{i+1}\right)$, then there are 2 adjacent vertices that have the same color, this contradicts the definition of vertex coloring. If

$$
l\left(x_{i}\right)=1 \quad i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \quad l\left(x_{i}\right)=2, \quad i \equiv 0(\bmod 2), \quad l\left(y_{i j}\right)=l\left(x_{i j}\right)=1,
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right), w\left(x_{i 1}\right) \neq w\left(x_{i 2}\right), w\left(x_{i 1}\right) \neq w\left(y_{i 2}\right)$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 7$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)= \begin{cases}1, & \text { for } i \equiv 1,3(\bmod 4) \text { or for } i \equiv 2(\bmod 4), \\
2, & \text { for } i \equiv 0(\bmod 4), \\
l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1 .\end{cases}
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}7, & \text { for } i=1, m, \\
8, & \text { for } i \equiv 0(\bmod 2), i \neq m, \\
9, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1,\end{cases} \\
w\left(x_{i j}\right)=\left\{\begin{array}{lr}
3, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=1,3 \text { or for } i \equiv 2(\bmod 4) \text { and } j=1,3, \\
4, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=2 \text { or for } i \equiv 2(\bmod 4) \text { and } j=2 \text { or } \\
5, & \text { for } i \equiv 0(\bmod 4) \text { and } j=1,3,
\end{array}\right. \\
w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=2, \\
3, & \text { for } i=0(\bmod 4) \text { and } 1 \leq j \leq 3 \text { or for } i \equiv 2(\bmod 4) \text { and } 1 \leq j \leq 3,\end{cases}
\end{gathered}
$$

We have the upper bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$ for $m \geq 4$ and $n=3$.
Case 8: $m=2$ and $n \geq 4$.
First step here is to find the lower bound of $V\left(P_{m} \bigodot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)<7$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$, if

$$
l\left(x_{1}\right)=1, \quad l\left(x_{2}\right)=2, \quad l\left(x_{i j}\right)=l\left(y_{i j}\right)=1,
$$

then $w\left(x_{i j+1}\right)=w\left(x_{i j+2}\right)$, then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{1}\right)=1, \quad l\left(x_{2}\right)=2, \quad l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \quad j \equiv 0(\bmod 2), \quad j=1, n, \\
l\left(y_{i j}\right)=2, \quad j \equiv 1,3(\bmod 4), \quad j \neq 1, \quad n \rightarrow w\left(x_{1}\right) \neq w\left(x_{2}\right), \quad w\left(x_{i j+1}\right) \neq w\left(x_{i j+2}\right),
\end{gathered}
$$

then we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 7$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)= \begin{cases}1, & \text { for } i=1, \\
2, & \text { for } i=2,\end{cases} \\
l\left(x_{i j}\right)=1,
\end{gathered} l\left(y_{i j}\right)=\left\{\begin{array}{l}
1, \text { for } i=1,2 \text { and } j=1, n \text { or for } i=1,2 \text { and } j \equiv 0(\bmod 2), \\
2, \text { for } i=1,2 \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n .
\end{array}\right.
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}3 n+1-n / 2, & \text { for } i=1 \text { and } n \equiv 0(\bmod 2), \\
3 n+1-\lceil n / 2\rceil, & \text { for } i=1 \text { and } n \equiv 1,3(\bmod 4), \\
2 n+n / 2, & \text { for } i=2 \text { and } n \equiv 0(\bmod 2), \\
2 n+\lfloor n / 2\rfloor, & \text { for } i=2 \text { and } n \equiv 1,3(\bmod 4),\end{cases} \\
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i=1 \text { and } j=1, n, \\
4, & \text { for } i=1 \text { and } j \equiv 0(\bmod 2), j \neq n \text { or for } i=2 \text { and } j=1, n, \\
5, & \text { for } i=1 \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n \text { or for } i=2 \text { and } j \equiv 0(\bmod 2), j \neq n, \\
6, & \text { for } i=2 \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n,\end{cases} \\
w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i=1 \text { and } 1 \leq j \leq n, \\
3, & \text { for } i=2 \text { and } 1 \leq j \leq n .\end{cases}
\end{gathered}
$$

The upper bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 7$ is true. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$ for $m=2$ and $n \geq 4$.
Case 9: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n=2$.
First step to prove this theorem in this case is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{\text {lis }}\left(P_{m} \odot C p_{n}\right)<7$, and let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$, if

$$
l\left(x_{i}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \quad l\left(x_{i}\right)=2, \quad i \equiv 3(\bmod 4), \quad l\left(x_{i j}\right)=l\left(y_{i j}\right)=1,
$$

then $w\left(x_{i 1}\right)=w\left(x_{i 2}\right), w\left(x_{i+1}\right)=w\left(x_{i+2}\right)$, then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \quad l\left(x_{i}\right)=2, \quad i \equiv 3(\bmod 4) \\
l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \quad 1 \leq i \leq m, \quad j=1, \quad l\left(y_{i j}\right)=2, \quad 1 \leq i \leq m, \quad j=2,
\end{gathered}
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right), w\left(x_{i j+1}\right) \neq w\left(x_{i j+2}\right)$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 7$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{aligned}
& l\left(x_{i}\right)= \begin{cases}1, & \text { for } i \equiv 1(\bmod 4) \text { or for } i \equiv 0(\bmod 2), \\
2, & \text { for } i \equiv 3(\bmod 4),\end{cases} \\
& w\left(x_{i j}\right)=1 ; \quad l\left(y_{i j}\right)= \begin{cases}1, & \text { for } 1 \leq i \leq m \text { and } j=1, \\
2, & \text { for } 1 \leq i \leq m \text { and } j=2 .\end{cases}
\end{aligned}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
w\left(x_{i}\right)= \begin{cases}6, & \text { for } i=1, m, \\ 7, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, \\ 8, & \text { for } i \equiv 0(\bmod 2), i \neq m,\end{cases}
$$

$$
\begin{aligned}
& w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i \equiv 1(\bmod 4) \text { and } j=1 \text { or for } i \equiv 0(\bmod 2) \text { and } j=1, \\
4, & \text { for } i \equiv 1(\bmod 4) \text { and } j=2 \text { or for } i \equiv 0(\bmod 2) \text { and } j=2 \text { or } \\
5, & \text { for } i \equiv 3(\bmod 4) \text { and } j=1,\end{cases} \\
& w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i \equiv 1(\bmod 4) \text { and } j=2, \\
3, & \text { for } i=3(\bmod 4) \text { and } j=1,2 \text { or for } i \equiv 0(\bmod 2) \text { and } j=1,2,\end{cases}
\end{aligned}
$$

The upper bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 7$ is true. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$ for $m \equiv 1,3(\bmod 4)$, $m \geq 5$ and $n=2$.

Case 10: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n=3$.
First step here is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)<7$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=6$, if

$$
l\left(x_{i}\right)=l\left(x_{i j}\right)=1, \quad l\left(y_{i 1}\right)=1, \quad l\left(y_{i j}\right)=2, \quad j=2,3,
$$

then $w\left(x_{i+1}\right)=w\left(x_{i+2}\right)$, then we have that there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{i}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \quad l\left(x_{i}\right)=2, \quad i \equiv 3(\bmod 4), \quad l\left(x_{i j}\right)=1 ; l\left(y_{i j}\right)=1,
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right), w\left(x_{i j+1}\right) \neq w\left(x_{i j+2}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 7$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=\left\{\begin{array}{ll}
1, & \text { for } i \equiv 1(\bmod 4) \text { or for } i \equiv 0(\bmod 2), \\
2, & \text { for } i \equiv 3(\bmod 4), \\
l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1 .
\end{array} .\right.
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
w\left(x_{i}\right)=\left\{\begin{array}{lll}
7, & & \text { for } i=1, m, \\
8, & & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, m, \\
9, & & \text { for } i \equiv 0(\bmod 2)
\end{array}\right.
$$

$$
\begin{gathered}
w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i \equiv 1(\bmod 4) \text { and } j=1,3 \text { or for } i \equiv 0(\bmod 2) \text { and } j=1,3, \\
4, & \text { for } i \equiv 1(\bmod 4) \text { and } j=2 \text { or for } i \equiv 0(\bmod 2) \text { and } j=2 \text { or } \\
\text { for } i \equiv 3(\bmod 4) \text { and } j=1,3,\end{cases} \\
5, \\
w\left(y_{i j}\right)= \begin{cases}2, & \text { for } i \equiv 3(\bmod 4) \text { and } j=2, \\
3, & \text { for } i=3(\bmod 4) \text { and } 1 \leq j \leq 3 \text { or for } i \equiv 0(\bmod 2) \text { and } 1 \leq j \leq 3,\end{cases} \\
w \leq 3 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 7 . \quad$ So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$ for $m \equiv 1,3(\bmod 4), m \geq 5$ and $n=3$.

Case 11: $m \equiv 0(\bmod 2) m \geq 4$ and $n \geq 4$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)<8$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$, if

$$
l\left(x_{i}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \quad l\left(x_{i}\right)=2, \quad i \equiv 0(\bmod 4), \quad l\left(x_{i j}\right)=l\left(y_{i j}\right)=1,
$$

then $w\left(x_{i j+1}\right)=w\left(x_{i j+2}\right)$, so there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \quad l\left(x_{i}\right)=2, \quad i \equiv 0(\bmod 4), \quad l\left(x_{i j}\right)=1, \\
l\left(y_{i j}\right)=1, \quad 1 \leq i \leq m, \quad j=1, n, \quad j \equiv 0(\bmod 2), \quad l\left(y_{i j}\right)=2, \\
1 \leq i \leq m, \quad j \equiv 1,3(\bmod 4), \quad j \neq 1, n,
\end{gathered}
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right), w\left(x_{i j+1}\right) \neq w\left(x_{i j+2}\right), w\left(x_{i j}\right) \neq w\left(y_{i j}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 8$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=\left\{\begin{array}{ll}
1, & \text { for } i \equiv 1,3(\bmod 4) \text { and } 1 \leq j \leq n \text { or for } i \equiv 2(\bmod 4) \text { and } 1 \leq j \leq n, \\
2, & \text { for } i \equiv 0(\bmod 4) \text { and } 1 \leq j \leq n,
\end{array} \quad l\left(x_{i j}\right)=1,\right.
\end{gathered} \quad l\left(y_{i j}\right)= \begin{cases}1, & \text { for } 1 \leq i \leq m \text { and } j=1, n \text { or for } 1 \leq i \leq m \text { and } j \equiv 0(\bmod 2), \\
2, & \text { for } 1 \leq i \leq m \text { and } j \equiv 1,3(\bmod 4), j \neq 1, m .\end{cases}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\begin{array}{c} 
\begin{cases}2 n+n / 2, & \text { for } i=1, m \text { and } n \equiv 0(\bmod 2), \\
2 n+\lfloor n / 2\rfloor, & \text { for } i=1, m \operatorname{and} n \equiv 1,3(\bmod 4), \\
3 n+1-n / 2, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1 \operatorname{and} n \equiv 0(\bmod 2), \\
3 n+1-\lceil n / 2\rceil, & \text { for } i \equiv 0(\bmod 2), i \neq m \text { and } n \equiv 1,3(\bmod 4), \\
3 n+2-n / 2, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1 \text { and } n \equiv 0(\bmod 2), \\
3 n+1-\lfloor n / 2\rfloor, & \text { for } i \equiv 0(\bmod 2), i \neq m \text { and } n \equiv 1,3(\bmod 4),\end{cases} \\
w\left(x_{i}\right)= \begin{cases}3, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=1, n \text { or for } i \equiv 2(\bmod 4) \text { and } j=1, n, \\
4, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j \equiv 0(\bmod 2), j \neq n \text { or }\end{cases} \\
\quad \text { for } i \equiv 2(\bmod 4) \text { and } j \equiv 0(\bmod 2), j \neq n \text { or for } i \equiv 0(\bmod 4) \text { and } j=1, n, \\
5, \\
\text { for } i \equiv 1,3(\bmod 4) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n \text { or } \\
\text { for } i \equiv 2(\bmod 4) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n \text { or } \\
\text { for } i \equiv 0(\bmod 4) \text { and } j \equiv 0(\bmod 2), j \neq n,
\end{array}, \quad \begin{array}{l}
\text { for } i \equiv 0(\bmod 4) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n,
\end{array}\right\} \begin{array}{ll}
2, & \text { for } i \equiv 1,3(\bmod 4) \text { and } 1 \leq j \leq n \text { or for } i \equiv 2(\bmod 4) \text { and } 1 \leq j \leq n, \\
3, & \text { for } i \equiv 0(\bmod 4) \text { and } 1 \leq j \leq n .
\end{array}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 8$. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=8$ for $m \equiv 0(\bmod 4), m \geq 4$ and $n \geq 4$.

Case 12: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n \geq 4$.
First step to prove this theorem in this case is to find the lower bound of $V\left(P_{m} \odot C P_{n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq \chi\left(P_{m} \odot C p_{n}\right)=3$.

Assume $\chi_{\text {lis }}\left(P_{m} \odot C p_{n}\right)<8$, let $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=7$, if

$$
l\left(x_{i}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \quad l\left(x_{i}\right)=2, \quad i \equiv 3(\bmod 4), \quad l\left(x_{i j}\right)=l\left(y_{i j}\right)=1,
$$

then $w\left(x_{i j+1}\right)=w\left(x_{i j+2}\right)$, so there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \quad l\left(x_{i}\right)=2, \quad i \equiv 3(\bmod 4), \quad l\left(x_{i j}\right)=1, \quad l\left(y_{i j}\right)=1, \\
1 \leq i \leq m, \quad j=1, n, \quad j \equiv 0(\bmod 2), \quad l\left(y_{i j}\right)=2, \quad 1 \leq i \leq m, \quad j \equiv 1,3(\bmod 4), \\
j \neq 1, n \rightarrow w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right), \quad w\left(x_{i j+1}\right) \neq w\left(x_{i j+2}\right), \quad w\left(x_{i j}\right) \neq w\left(y_{i j}\right),
\end{gathered}
$$

therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \geq 8$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot C p_{n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=\left\{\begin{array}{ll}
1, & \text { for } i \equiv 1(\bmod 4) \text { or for } i \equiv 0(\bmod 2), \\
2, & \text { for } i \equiv 3(\bmod 4),
\end{array} \quad l\left(x_{i j}\right)=1,\right.
\end{gathered} l\left(y_{i j}\right)= \begin{cases}1, & \text { for } 1 \leq i \leq m \text { and } j=1, n \text { or for } 1 \leq i \leq m \text { and } j \equiv 0(\bmod 2), \\
2, & \text { for } 1 \leq i \leq m \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n .\end{cases}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{aligned}
& w\left(x_{i}\right)= \begin{cases}2 n+n / 2, & \text { for } i=1, m \text { and } n \equiv 0(\bmod 2), \\
3 n+2-n / 2, & \text { for } i=0(\bmod 2) \text { and } n \equiv 0(\bmod 2), \\
3 n+1-n / 2, & \text { for } i \equiv 1,3(\bmod 4) \text { and } n \equiv 0(\bmod 2), \\
2 n+\lfloor n / 2\rfloor, & \text { for } i=1, m \operatorname{and} n \equiv 1,3(\bmod 4), \\
3 n-\lfloor n / 2\rfloor, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1 \operatorname{and} n \equiv 1,3(\bmod 4), \\
3 n+1-\lfloor n / 2\rfloor, & \text { for } i=0(\bmod 4) \text { and } n \equiv 1,3(\bmod 4),\end{cases} \\
& w\left(x_{i j}\right)= \begin{cases}3, & \text { for } i \equiv 1(\bmod 4) \text { and } j=1, n \text { or for } i \equiv 0(\bmod 2) \text { and } j=1, n, \\
4, & \text { for } i \equiv 1(\bmod 4) \text { and } j \equiv 0(\bmod 2), j \neq n \text { or } \\
\text { for } i \equiv 0(\bmod 2) \text { and } j \equiv 0(\bmod 2), j \neq n \text { or } \\
5, & \text { for } i \equiv 1(\bmod 4) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n \text { or }\end{cases} \\
& w\left(y_{i j}\right)=\left\{\begin{array}{rr}
2, & \text { for } i \equiv 0(\bmod 2) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n \text { or } \\
6, & \text { for } i \equiv 3(\bmod 4) \text { and } j \equiv 0(\bmod 2), j \neq n,
\end{array}\right. \\
& 3,3(\bmod 4) \text { and } j \equiv 1,3(\bmod 4), j \neq 1, n,
\end{aligned}, \begin{aligned}
& \text { for } 4) \text { and } 1 \leq j \leq n .
\end{aligned}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot C p_{n}\right) \leq 8$. So $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)=8$ for $m \geq 5$ and $n \geq 4$.

Theorem 2. Let $G=P_{m} \odot S_{2, n}$ for $n, m \geq 2$, then the chromatic number of local irregular $G$ is

$$
\chi_{l i s}\left(P_{m} \bigodot S_{2, n}\right)= \begin{cases}5, & \text { for } m=3 \text { and } n \geq 2 \\ 6, & \text { for } m=2 \text { and } n \geq 2, \\ 7, & \text { for } m \geq 4 \text { and } n \geq 2,\end{cases}
$$

with $\operatorname{opt}(l)\left(P_{m} \odot S_{2, n}\right)=1,2$, for $m \geq 2$ and $n \geq 2$.
Proof. Vertex set is

$$
\begin{gathered}
V\left(P_{n} \bigodot S_{2, n}\right)=\left\{x_{i} ; 1 \leq i \leq m\right\} \cup\left\{a_{i} ; 1 \leq i \leq m\right\} \cup\left\{b_{i} ; 1 \leq i \leq m\right\} \\
\cup\left\{a_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{b_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\}
\end{gathered}
$$

and the edge set is

$$
\begin{gathered}
E\left(P_{n} \bigodot S_{2, n}\right)=\left\{x_{i} x_{i+1}, 1 \leq i \leq m-1\right\} \cup\left\{a_{i} b_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i} a_{i} ; 1 \leq i \leq m\right\} \\
\cup\left\{x_{i} b_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i} a_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{x_{i} b_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \\
\cup\left\{a_{i} a_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{b_{i} b_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} .
\end{gathered}
$$

The order and the size respectively are $2 m n+3 m$ and $4 m n+4 m-1$. This proof is divided into 4 cases as follows.

Case 1: $m=3$ and $n \geq 2$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot S_{2, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq \chi\left(P_{m} \odot S_{2, n}\right)=3$.

Assume $\chi l i s\left(P_{m} \odot S_{2, n}\right)=4$, if $l\left(a_{i}\right)=l\left(b_{i}\right)=1, l\left(x_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=1$ then $w\left(a_{i}\right)=w\left(b_{i}\right)$, then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad 1 \leq j \leq n-1, \quad l\left(b_{i n}\right)=2,
$$

then

$$
w\left(a_{i}\right) \neq w\left(b_{i}\right), \quad w\left(x_{1}\right)=w\left(x_{3}\right) \neq w\left(x_{2}\right),
$$

therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq 5$.
After that, we will find the upper bound of $\chi l i s\left(P_{m} \odot S_{2, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{2, n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{aligned}
& l\left(x_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \\
& l\left(b_{i j}\right)= \begin{cases}1, & \text { for } 1 \leq i \leq 3 \text { and } 1 \leq j \leq n-1, \\
2, & \text { for } 1 \leq i \leq 3 \text { and } j=n .\end{cases}
\end{aligned}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}2 n+4, & \text { for } i=1,3, \\
2 n+5, & \text { for } i=2,\end{cases} \\
w\left(a_{i}\right)=n+2, \quad \text { for } 1 \leq i \leq 3, \\
w\left(b_{i}\right)=n+3, \quad \text { for } 1 \leq i \leq 3, \\
w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2 .
\end{gathered}
$$

The upper bound $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \leq 5$ is true. So $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=5$ for $m=3$ and $n \geq 2$.
Case 2: $m=2$ and $n \geq 2$.
First step here is to find the lower bound of $V\left(P_{m} \odot S_{2, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq \chi\left(P_{m} \odot S_{2, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=5$, if

$$
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=l\left(b_{2 j}\right)=1, \quad l\left(b_{1 j}\right)=1, \quad 1 \leq j \leq n-1, \quad l\left(b_{1 n}\right)=2,
$$

and then $w\left(a_{2}\right)=w\left(b_{2}\right)$, and there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{aligned}
& l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=1, \\
& l\left(b_{1, j}\right)=1, \quad l\left(b_{1, n}\right)=2, \quad l\left(b_{2, j}\right)=2, \quad j=1, n, \quad l\left(b_{2 j}\right)=1, \quad 2 \leq j \leq n-1,
\end{aligned}
$$

then $w\left(a_{i}\right) \neq w\left(b_{i}\right), w\left(x_{1}\right) \neq w\left(x_{2}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq 6$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{2, n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \\
l\left(b_{i j}\right)= \begin{cases}1, & \text { for } i=1 \text { and } 1 \leq j \leq n-1 \text { or for } i=2 \text { and } 2 \leq j \leq n-1, \\
2, & \text { for } i=1 \text { and } j=n \text { or for } i=2 \text { and } j=1, n .\end{cases}
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}2 n+4, & \text { for } i=1, \\
2 n+5, & \text { for } i=2,\end{cases} \\
w\left(a_{i}\right)={ }^{n+2,} \text { for } i=1,2, \\
w\left(b_{i}\right)= \begin{cases}n+3, & \text { for } i=1, \\
n+4, & \text { for } i=2,\end{cases} \\
w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \leq 6$. So $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=6$ for $m=2$ and $n \geq 2$.
Case 3: $m \equiv 0(\bmod 4), m \geq 4$ and $n \geq 2$.
First step to prove this theorem in this case is to find the lower bound of $V\left(P_{m} \odot S_{2, n}\right)$. Based on Lemma 1, we have $\chi_{\text {lis }}\left(P_{m} \odot S_{2, n}\right) \geq \chi\left(P_{m} \odot S_{2, n}\right)=3$.

Assume $\chi_{\text {lis }}\left(P_{m} \odot S_{2, n}\right)=6$, if

$$
\begin{gathered}
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad i \equiv 2(\bmod 4), \\
l\left(b_{i j}\right)=1, \quad i \equiv 0(\bmod 4), \quad j \neq 1, n, \quad l\left(b_{i j}\right)=2, \quad i \equiv 0(\bmod 4), \quad j=1, n,
\end{gathered}
$$

then $w\left(a_{i}\right)=w\left(b_{i}\right)$, so there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad i=m, \quad 1 \leq j \leq n-1, \\
i \equiv 0(\bmod 2), \quad i \neq m, \quad 2 \leq j \leq n-1, \quad l\left(b_{i j}\right)=2, \quad i \equiv 1,3(\bmod 4), \\
i=m, \quad j=n, \quad i \equiv 0(\bmod 2), \quad i \neq m, \quad j=1, n,
\end{gathered}
$$

then $w\left(a_{i}\right) \neq w\left(b_{i}\right), w\left(x_{i}\right) \neq w\left(x_{i+1}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq 7$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)$
Furthermore, we define $l: V\left(P_{m} \odot S_{2, n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \\
l\left(b_{i j}\right)=\left\{\begin{array}{l}
1, \quad \text { for } i \equiv 1,3(\bmod 4) \text { and } 1 \leq j \leq n-1 \text { or } \\
\quad \text { for } i \equiv 0(\bmod 2), i \neq m \text { and } 2 \leq j \leq n-1 \text { or } \\
2, \quad \text { for } i=m, \text { and } 1 \leq j \leq n-1, \\
\text { for } i \equiv 0(\bmod 4) \text { and } j=n \text { or } \\
\text { for } i=m, \text { and } j=n .
\end{array}\right.
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\begin{array}{c}
w\left(x_{i}\right)= \begin{cases}2 n+4, & \text { for } i=1, m, \\
2 n+5, & \text { for } i \equiv 0(\bmod 2), i \neq m, \\
2 n+6, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1,\end{cases} \\
w\left(a_{i}\right)=n+2, \quad \text { for } 1 \leq i \leq m,
\end{array}\right\}\left(b_{i}\right)=\left\{\begin{array}{l}
n+3, \quad \text { for } i \equiv 1,3(\bmod 4), i=m, \\
n+4, \\
\text { for } i \equiv 0(\bmod 2), i \neq m,
\end{array}\right\}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=7$ for $m \equiv 0(\bmod 2)$, $m \geq 4$ and $n \geq 2$.

Case 4: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n \geq 2$.
First step here is to find the lower bound of $V\left(P_{m} \bigodot S_{2, n}\right)$. Based on Lemma 1, we have

$$
\chi_{l i s}\left(P_{m} \bigodot S_{2, n}\right) \geq \chi\left(P_{m} \bigodot S_{2, n}\right)=3
$$

Assume $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=6$, if

$$
\begin{aligned}
& l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \quad i \equiv 1(\bmod 4), \quad i \equiv 0(\bmod 2), \\
& \quad l\left(b_{i j}\right)=1, \quad i \equiv 3(\bmod 4), \quad j \neq n, \quad l\left(b_{i j}\right)=2, \quad i \equiv 3(\bmod 4), \quad j=n,
\end{aligned}
$$

then $w\left(a_{i}\right)=w\left(b_{i}\right)$, and there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad 1 \leq j \leq n-1,
$$

$$
i \equiv 0(\bmod 2), \quad 2 \leq j \leq n-1, \quad l\left(b_{i j}\right)=2, \quad i \equiv 1,3(\bmod 4), \quad j=n, \quad i \equiv 0(\bmod 2), \quad j=1, n,
$$ then $w\left(a_{i}\right) \neq w\left(b_{i}\right), w\left(x_{i}\right) \neq w\left(x_{i+1}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \geq 7$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{2, n}\right) \rightarrow\{1,2\}$ with the vertex irregular 2-labelling as follows:

$$
l\left(x_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(a_{i j}\right)=1,
$$

$l\left(b_{i j}\right)= \begin{cases}1, & \text { for } i \equiv 1,3(\bmod 4) \text { and } 1 \leq j \leq n-1 \text { or for } i \equiv 0(\bmod 2), \text { and } 2 \leq j \leq n-1, \\ 2, & \text { for } i \equiv 1,3(\bmod 4) \text { and } j=n \text { or for } i \equiv 0(\bmod 2) \text { and } j=1, n .\end{cases}$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{aligned}
& w\left(x_{i}\right)= \begin{cases}2 n+4, & \text { for } i=1, m, \\
2 n+5, & \text { for } i \equiv 1,3(\bmod 4), \\
2 n+6, & \text { for } i \equiv 0(\bmod 2),\end{cases} \\
& w\left(a_{i}\right)=n+2, \quad \text { for } 1 \leq i \leq m, \\
& w\left(b_{i}\right)= \begin{cases}n+3, & \text { for } i \equiv 1,3(\bmod 4), \\
n+4, & \text { for } i \equiv 0(\bmod 2),\end{cases} \\
& w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2 .
\end{aligned}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=7$ for $m \equiv 1,3(\bmod 4)$, $m \geq 5$ and $n \geq 2$.

Theorem 3. Let $G=P_{m} \odot S_{3, n}$ for $n, m \geq 2$, then the chromatic number of local irregular $G$ is

$$
\chi_{l i s}\left(P_{m} \bigodot S_{3, n}\right)= \begin{cases}5, & \text { for } m=3 \text { and } n \geq 2 \\ 6, & \text { for } m=2 \text { and } n \geq 2 \\ 7, & \text { for } m \geq 4 \text { and } n \geq 3\end{cases}
$$

with

$$
\operatorname{opt}(l)\left(P_{m} \bigodot S_{3, n}\right)= \begin{cases}1, & \text { for } m=3 \text { and } n=3, \\ 1,2, & \text { for } m=2 \text { and } n=2 \text { or for } m=3 \text { and } n=2 \text { or } \\ & \text { for } m \geq 4 \text { and } n \geq 2 .\end{cases}
$$

Proof. The vertex set is

$$
\begin{aligned}
& V\left(P_{m} \bigodot S_{3, n}\right)=\left\{x_{i} ; 1 \leq i \leq m\right\} \cup\left\{a_{i} ; 1 \leq i \leq m\right\} \cup\left\{b_{i} ; 1 \leq i \leq m\right\} \cup\left\{c_{i} ; 1 \leq i \leq m\right\} \\
& \cup\left\{a_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{b_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\} \cup\left\{c_{i j} ; 1 \leq i \leq m, 1 \leq j \leq n\right\}
\end{aligned}
$$

and the edge set is

$$
\begin{gathered}
V\left(P_{m} \bigodot S_{3, n}\right)=\left\{x_{i} x_{i+1} ; 1 \leq i \leq m-1\right\} \cup\left\{x_{i} y_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i} a_{i} ; 1 \leq i \leq m\right\} \\
\cup\left\{x_{i} b_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i} c_{i} ; 1 \leq i \leq m\right\} \cup\left\{y_{i} a_{i} ; 1 \leq i \leq m\right\} \cup\left\{y_{i} b_{i} ; 1 \leq i \leq m\right\} \\
\cup\left\{y_{i} c_{i} ; 1 \leq i \leq m\right\} \cup\left\{x_{i} a_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} \cup\left\{x_{i} b_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} \\
\cup\left\{x_{i} c_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} \cup\left\{a_{i} a_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} \\
\cup\left\{b_{i} b_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} \cup\left\{c_{i} c_{i j} ; 1 \leq i \leq m ; 1 \leq j \leq n\right\} .
\end{gathered}
$$

The order and size respectively are $3 m n+5 m$ and $6 m n+8 n-1$. This proof can be divided into 8 following cases.

Case 1: $m=3$ and $n=2$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=4$, if

$$
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(y_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(c_{i j}\right)=1,
$$

then $w\left(a_{i}\right)=w\left(b_{i}\right)=w\left(c_{i}\right)=w\left(y_{i}\right)$, and there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(c_{i j}\right)=1, \quad l\left(y_{i}\right)=2,
$$

then $\left(w\left(a_{i}\right)=w\left(b_{i}\right)=w\left(c_{i}\right)\right) \neq w\left(y_{i}\right), w\left(x_{1}\right) \neq w\left(x_{2}\right)$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 5$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=2, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1 .
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)=\left\{\begin{array}{l}
12, \text { for } i=1,3, \\
13, \\
\text { for } i=2,
\end{array}\right. \\
w\left(y_{i}\right)=4, \quad w\left(a_{i}\right)=5, \quad w\left(b_{i}\right)=5, \quad w\left(c_{i}\right)=5, \quad w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 5$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=5$ for $m=3$ and $n=2$.
Case 2: $m=3$ and $n=3$.
Based on Proposition 1, we have $\operatorname{opt}(l)=1$. So the lower bound ( $P_{m} \odot S_{3, n}$ ) is $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 5$
Since $\operatorname{opt}(l)=1$, the labelling provides the vertex-weight as follows:

$$
\begin{gathered}
w\left(x_{i}\right)= \begin{cases}3 n+5, & \text { for } i=1,3, \\
3 n+6, & \text { for } i=2,\end{cases} \\
w\left(y_{i}\right)=4,
\end{gathered} \begin{gathered}
w\left(a_{i}\right)=n+1 \text { for } 1 \leq i \leq 3, \\
w\left(b_{i}\right)=n+1 \text { for } 1 \leq i \leq 3, \\
w\left(c_{i}\right)=n+1 \text { for } 1 \leq i \leq 3, \\
w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 5$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=5$ for $m=3$ and $n \geq 2$.
Case 3: $m=2$ and $n=2$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=5$, if

$$
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(c_{i j}\right)=1, \quad l\left(y_{1}\right)=1, \quad l\left(y_{2}\right)=2,
$$

then $w\left(a_{2}\right)=w\left(b_{2}\right)=w\left(c_{2}\right)=w\left(y_{2}\right)$ and there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(b_{i j}\right)=1, \quad l\left(y_{i}\right)=2, \quad l\left(c_{1 j}\right)=1, \quad l\left(c_{2,1}\right)=1, \quad l\left(c_{2,2}\right)=2,
$$

then $w\left(x_{1}\right) \neq w\left(x_{2}\right), w\left(y_{i}\right) \neq\left(\left(w\left(a_{i}\right)=w\left(b_{i}\right)=w\left(c_{i}\right)\right)\right.$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 6$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.

Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
\begin{array}{cc}
l\left(x_{i}\right)=1, & l\left(y_{i}\right)=2, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \\
l\left(c_{i j}\right)= \begin{cases}1, & \text { for } i=1 \text { and } j=1,2 \text { or for } i=2 \text { and } j=1, \\
2, & \text { for } i=2 \text { and } j=2 .\end{cases}
\end{array}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\begin{aligned}
& w\left(x_{i}\right)= \begin{cases}12, & \text { for } i=1, \\
13, & \text { for } i=2,\end{cases} \\
& w\left(y_{i}\right)=4, \quad w\left(a_{i}\right)=5, \quad w\left(b_{i}\right)=5, \\
& w\left(c_{i}\right)= \begin{cases}5, & \text { for } i=1, \\
6, & \text { for } i=2,\end{cases} \\
& w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2 .
\end{aligned}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right) \leq 6$. So $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)=6$ for $m=2$ and $n=2$.
Case 4: $m=2$ and $n \geq 3$.
First step here is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=5$, if
$l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(y_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i=1,2, \quad 1 \leq j \leq n-1$ $l\left(c_{i j}\right)=2, \quad i=1,2, \quad j=n$,
then $w\left(x_{1}\right)=w\left(x_{2}\right)$, then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{array}{ll}
l\left(x_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i=1, \quad 1 \leq j \leq n, \quad i=2, \quad 1 \leq j \leq n-1, \\
& l\left(c_{i j}\right)=2, \quad i=2, \quad j=n,
\end{array}
$$

then $w\left(x_{1}\right) \neq w\left(x_{2}\right), w\left(y_{i}\right) \neq\left(\left(w\left(a_{i}\right)=w\left(b_{i}\right)=w\left(c_{i}\right)\right)\right.$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 6$.

After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
\begin{aligned}
& l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \\
& \quad l\left(c_{i j}\right)= \begin{cases}1, & \text { for } i=1 \text { and } 1 \leq j \leq n \text { or for } i=2 \text { and } 1 \leq j \leq n-1, \\
2, & \text { for } i=2 \text { and } j=n .\end{cases}
\end{aligned}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\begin{array}{c}
w\left(x_{i}\right)= \begin{cases}3 n+5, & \text { for } i=1, \\
3 n+6, & \text { for } i=2,\end{cases} \\
w\left(a_{i}\right)=n+1, \quad \text { for } i=1,2,
\end{array}, \begin{array}{c}
w\left(y_{i}\right)=4, \\
w\left(b_{i}\right)=n+1, \quad \text { for } i=1,2,
\end{array}\right\} \begin{array}{ll}
n+1, & \text { for } i=1, \\
n+2, & \text { for } i=2,
\end{array}, ~ \begin{gathered}
\\
w\left(c_{i}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2 .
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 6$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=6$ for $m=2$ and $n \geq 3$.
Case 5: $m \equiv 0(\bmod 2) m \geq 4$ and $n=2$.
First step to prove this theorem in this case is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=6$, if

$$
\begin{gathered}
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(y_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \\
j=1,2, \quad i \equiv 0(\bmod 4), \quad j=1,2, \quad l\left(c_{i j}\right)=1, \quad i \equiv 2(\bmod 4), \\
j=1, \quad l\left(c_{i j}\right)=2, \quad i \equiv 2(\bmod 4), \quad j=2,
\end{gathered}
$$

then $w\left(y_{i}\right)=w\left(a_{i}\right)$. Then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=1, \quad l\left(y_{i}\right)=2, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \\
i \equiv 0(\bmod 2), \quad j=1, \quad i \neq m, \quad i \equiv 1,3(\bmod 4), \quad j=1,2, \quad l\left(c_{i j}\right)=2, \\
i \equiv 0(\bmod 2), \quad i \neq m, \quad j=2,
\end{gathered}
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2} ; w\left(y_{i}\right) \neq w\left(a_{i}\right)\right.$. Therefore we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 7$. After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.

Furthermore, we define $l: V\left(P_{m} \bigodot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
\begin{aligned}
& l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=2, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \\
& l\left(c_{i j}\right)=\left\{\begin{array}{l}
1, \quad \text { for } i \equiv 1,3(\bmod 4) \text { and } j=1,2 \text { or for } i=m \text { and } j=1,2 \text { or } \\
2, \quad \text { for } i \equiv 0(\bmod 2), \quad i \neq m \text { and } j=1,
\end{array}\right.
\end{aligned}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\begin{array}{c}
w\left(x_{i}\right)= \begin{cases}12, & \text { for } i=1, m, \\
13, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, \\
14, & \text { for } i \equiv 0(\bmod 2), i \neq m,\end{cases} \\
w\left(y_{i}\right)=4, \quad w\left(a_{i}\right)=5, \quad w\left(b_{i}\right)=5,
\end{array}\right\} \begin{array}{ll}
5, & \text { for } i \equiv 1,3(\bmod 4), \\
6, & \text { for } i \equiv 0(\bmod 2), i \neq m,
\end{array}, \begin{gathered}
\\
w\left(c_{i}\right)=\left\{\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2,\right.
\end{gathered}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=7$ for $m \equiv 0(\bmod 2) ; m \geq 4$ and $n=2$.

Case 6: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n=2$.
First step here is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=6$, if

$$
\begin{gathered}
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(y_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \\
i \equiv 1(\bmod 4), \quad j=1,2, \quad i \equiv 0(\bmod 2), \quad j=1,2, \quad l\left(c_{i j}\right)=1, \\
i \equiv 3(\bmod 4), \quad j=1, \quad l\left(c_{i j}\right)=2, \quad i \equiv 3(\bmod 4), \quad j=2,
\end{gathered}
$$

then $w\left(y_{i}\right)=w\left(a_{i}\right)$, then there are 2 adjacent vertices that have same color, it contradicts to definition of vertex coloring. If

$$
\begin{aligned}
& l\left(x_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \\
& j=1,2, \quad i \equiv 0(\bmod 2), \quad j=1, \quad l\left(c_{i j}\right)=2, \quad i \equiv 0(\bmod 2), \quad j=2, \quad l\left(y_{i}\right)=2,
\end{aligned}
$$

then

$$
w\left(x_{i+1}\right) \neq w\left(x_{i+2}, \quad w\left(y_{i}\right) \neq w\left(a_{i}\right), \quad w\left(y_{i}\right) \neq w\left(b_{i}\right), \quad w\left(y_{i}\right) \neq w\left(c_{i}\right) .\right.
$$

We have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 7$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=2, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \\
l\left(c_{i j}\right)=\left\{\begin{array}{l}
1, \quad \text { for } i \equiv 1,3(\bmod 4) \text { and } j=1,2 \text { or } \\
\quad \text { for } i \equiv 0(\bmod 2), \quad i \neq m \text { and } j=1, \\
2, \quad \text { for } i \equiv 0(\bmod 2), \quad i \neq m \text { and } j=2 .
\end{array}\right.
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\begin{array}{c}
w\left(x_{i}\right)= \begin{cases}12, & \text { for } i=1, m \\
13, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, m, \\
14, & \text { for } i \equiv 0(\bmod 2),\end{cases} \\
w\left(y_{i}\right)=4, \quad w\left(a_{i}\right)=5, \quad w\left(b_{i}\right)=5,
\end{array}\right\} \begin{array}{ll}
w\left(c_{i}\right)= \begin{cases}5, & \text { for } i \equiv 1,3(\bmod 4), i=1, m, \\
6, & \text { for } i \equiv 0(\bmod 2),\end{cases} \\
w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2 .
\end{array}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 7 . \quad$ So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=7$ for $m \equiv 1,3(\bmod 4) ; m \geq 5$ and $n=2$.

Case 7: $m \equiv 0(\bmod 2) m \geq 4$ and $n \geq 3$.
First step to prove this theorem is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=6$, it is true if

$$
\begin{gathered}
l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=l\left(y_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=1, \quad l\left(c_{i j}\right)=1, \quad 1 \leq i \leq m, \\
1 \leq j \leq n-1, \quad l\left(c_{i j}\right)=2, \quad 1 \leq i \leq m, \quad j=n,
\end{gathered}
$$

then $w\left(x_{i+1}\right)=w\left(x_{i+2}\right)$, then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{array}{rc}
l\left(x_{i}\right)=1, & l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(y_{i}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \quad 1 \leq j \leq n, \\
i \equiv 0(\bmod 2), & i \neq m, \quad 1 \leq j \leq n-1, \quad i=m, \quad 1 \leq j \leq n, \quad l\left(c_{i j}\right)=2, \quad i \equiv 0(\bmod 2), \\
i \neq m, \quad i \neq m, \quad j=n, \quad w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right),
\end{array}
$$

we have the lower bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 7$. After that, we will find the upper bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.

Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1
$$

$$
l\left(c_{i j}\right)= \begin{cases}1, & \text { for } i \equiv 1,3(\bmod 4) \text { and } 1 \leq j \leq n \text { or } \\ \quad & \text { for } i \equiv 0(\bmod 2), i \neq m, \text { and } 1 \leq j \leq n-1 \text { or } \\ \quad & \text { for } i=m, \text { and } 1 \leq j \leq n, \\ 2, & \text { for } i \equiv 0(\bmod 2), i \neq m \text { and } j=n .\end{cases}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
\left.\left.\left.\begin{array}{c}
w\left(x_{i}\right)= \begin{cases}3 n+5, & \text { for } i=1, m, \\
3 n+6, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, \\
3 n+7, & \text { for } i \equiv 0(\bmod 2), i \neq m,\end{cases} \\
w\left(y_{i}\right)=4,
\end{array}\right\} \begin{array}{c}
w\left(a_{i}\right)=n+2, \quad \text { for } 1 \leq i \leq m,
\end{array}\right\} \begin{array}{l}
w\left(b_{i}\right)=n+2, \quad \text { for } 1 \leq i \leq m,
\end{array}\right\}\left(c_{i}\right)=\left\{\begin{array}{ll}
n+2, & \text { for } i=m, \text { or for } i \equiv 1,3(\bmod 4), \\
n+3, & \text { for } i \equiv 0(\bmod 2), i \neq m,
\end{array}\right\}
$$

The upper bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=7$ for $m \equiv 0(\bmod 2) ; m \geq 4$ and $n \geq 3$.

Case 8: $m \equiv 1,3(\bmod 4), m \geq 5$ and $n \geq 3$.
First step to prove the theorem in this case is to find the lower bound of $V\left(P_{m} \odot S_{3, n}\right)$. Based on Lemma 1, we have $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq \chi\left(P_{m} \odot S_{3, n}\right)=3$.

Assume $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=6$, if

$$
l\left(x_{i}\right)=l\left(y_{i}\right)=l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(c_{i j}\right)=1, \quad l\left(a_{i}\right)=l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=2,
$$

then $w\left(x_{i+1}\right)=w\left(x_{i+2}\right)$. Then there are 2 adjacent vertices that have same color, it contradicts the definition of vertex coloring. If

$$
\begin{gathered}
l\left(x_{i}\right)=l\left(a_{i}\right)=l\left(b_{i}\right)=l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=l\left(b_{i j}\right)=l\left(y_{i}\right)=1, \quad l\left(c_{i j}\right)=1, \quad i \equiv 1,3(\bmod 4), \\
1 \leq j \leq n, \quad i \equiv 0(\bmod 2), \quad 1 \leq j \leq n-1, \quad l\left(c_{i j}\right)=2, \quad i \equiv 0(\bmod 2), \quad j=n,
\end{gathered}
$$

then $w\left(x_{i+1}\right) \neq w\left(x_{i+2}\right)$. Based on that we have the lower bound $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \geq 7$.
After that, we will find the upper bound of $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.
Furthermore, we define $l: V\left(P_{m} \odot S_{3, n}\right) \rightarrow\{1,2\}$ with vertex irregular 2-labelling as follows:

$$
\begin{gathered}
l\left(x_{i}\right)=1, \quad l\left(y_{i}\right)=1, \quad l\left(a_{i}\right)=1, \quad l\left(b_{i}\right)=1, \quad l\left(c_{i}\right)=1, \quad l\left(a_{i j}\right)=1, \quad l\left(b_{i j}\right)=1, \\
l\left(c_{i j}\right)=\left\{\begin{array}{l}
1, \quad \text { for } i \equiv 1,3(\bmod 4), \text { and } 1 \leq j \leq n \text { or } \\
\quad \text { for } i \equiv 0(\bmod 2), \text { and } 1 \leq j \leq n-1, \\
2, \quad \text { for } i \equiv 0(\bmod 2), \quad \text { and } j=n .
\end{array}\right.
\end{gathered}
$$

Hence, $\operatorname{opt}(l)=2$ and the labelling provides the vertex-weight as follows:

$$
w\left(x_{i}\right)= \begin{cases}3 n+5, & \text { for } i=1, m, \\ 3 n+6, & \text { for } i \equiv 1,3(\bmod 4), i \neq 1, m, \\ 3 n+7, & \text { for } i \equiv 0(\bmod 2)\end{cases}
$$

$$
\left.\begin{array}{c}
w\left(y_{i}\right)=4 \\
w\left(a_{i}\right)=n+2, \text { for } 1 \leq i \leq m \\
w\left(b_{i}\right)=n+2, \text { for } 1 \leq i \leq m
\end{array}\right\} \begin{aligned}
& w\left(c_{i}\right)=\left\{\begin{array}{l}
n+2, \quad \text { for } i \equiv 1,3(\bmod 4) \\
n+3, \\
\text { for } i \equiv 0(\bmod 2)
\end{array}\right. \\
& w\left(a_{i j}\right)=2, \quad w\left(b_{i j}\right)=2, \quad w\left(c_{i j}\right)=2
\end{aligned}
$$

The upper bound is true: $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right) \leq 7$. So $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)=7$ for $m \equiv 1,3(\bmod 4)$, $m \geq 5$ and $n \geq 3$.

## 3. Conclusion

In this paper, we have studied the coloring of the vertices of the local irregular corona product by the graph of the family tree. We determined the exact value of the local irregular chromatic number of the corona product from the graph of the family tree, namely $\chi_{l i s}\left(P_{m} \odot C p_{n}\right)$, $\chi_{l i s}\left(P_{m} \odot S_{2, n}\right)$ and $\chi_{l i s}\left(P_{m} \odot S_{3, n}\right)$.

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